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# C++ Language Support for Generic Programming

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# What is Generic Programming?

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- Paradigm for software development
- Express algorithms (and data structures) in terms of abstract requirements
  - Applicable to many data types
  - Runtime performance of concrete implementations
  - Alternative implementations for special-case data structures
- Distinct from template metaprogramming





# Outline

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- Introduction to GP extensions for C++
  - Generic functions & types
  - Concepts
  - Models
- Existing practice & the Standard Library
- Impact on the C++ community



# Introduction to Generic Programming Extensions for C++

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# Vector equality

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- `std::vector` has this equality operator:

```
template<typename T, typename Alloc>
    bool operator==(const vector<T, Alloc>& x,
                    const vector<T, Alloc>& y);
```

- When can we instantiate this function?
  - Need to be able to compare `T` objects.
  - This requirement is **implicit** in the source code



# Explicit requirements

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- Make the requirement explicit:

```
template<typename T, typename Alloc>
    where { EqualityComparable<T> }
    bool operator==(const vector<T, Alloc>& x,
                    const vector<T, Alloc>& y);
```

- where clause states requirements
  - EqualityComparable is a **concept**
  - “T must be EqualityComparable”



# Equality Comparable concept

---

```
template<typeid T>
concept EqualityComparable
{
    bool operator==(const T&, const T&);
    bool operator!=(const T&, const T&);

    // == is an equivalence relation
    // != is the complement of ==
};
```

- A **concept** gives a name to a set of requirements
  - Syntax
  - Semantics
- A type **models** the concept if it satisfies the requirements



# Type-checking calls

- We can now type-check calls to :

```
template<typename T, typename Alloc>
  where { EqualityComparable<T> }
  bool operator==(const vector<T, Alloc>& x,
                  const vector<T, Alloc>& y);
```

```
vector<int> iv1, iv2;
if (iv1 == iv2) { ... }
vector<my_type> mv1, mv2;
if (mv1 == mv2) { ... }
```

- We need to ensure that `int` and `my_type` model `EqualityComparable`





# Models of EqualityComparable

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- `int models EqualityComparable`
  - The Standard Library asserts this
- Does `my_type` model `EqualityComparable`?
  - Only if syntax and semantics match
  - If yes, write a **model definition**:

```
model EqualityComparable<my_type> { };
```



# Opaque template parameters

- Let's add the implementation:

```
template<typeid T, typeid Alloc>
  where { EqualityComparable<T> }
  bool operator==(const vector<T, Alloc>& x,
                  const vector<T, Alloc>& y)
  {
    return x.size() == y.size()
           && equal(x.begin(), x.end(), y.begin());
  }
```

- `typename` eliminates type checking: you can do anything to `typename` types
- `typeid` enables type checking: you can do nothing but what you require
- Two-phase type checking: like C++ already has!



# What about the call to `equal`?

- `equal` has this signature with concepts:

```
template<typeid Iter1, typeid Iter2>
  where { EqualityComparable2<
           InputIterator<Iter1>::value_type,
           InputIterator<Iter2>::value_type> }
bool
equal(Iter1 first1, Iter1 last1, Iter2 first2);
```

- `EqualityComparable2` is a new concept:

```
template<typeid T, typeid U>
concept EqualityComparable2
{
  bool operator==(const T&, const U&);
  bool operator!=(const T&, const U&);
};
```



# What about the call to `equal`?

- `equal` has this signature with concepts:

```
template<typeid Iter1, typeid Iter2>
  where { EqualityComparable2<
           InputIterator<Iter1>::value_type,
           InputIterator<Iter2>::value_type> }
  bool
  equal(Iter1 first1, Iter1 last1, Iter2 first2);
```

- Both `Iter1` and `Iter2` must model the `InputIterator` concept.



# What about the call to `equal`?

- `equal` has this signature with concepts:

```
template<typeid Iter1, typeid Iter2>
  where { EqualityComparable2<
           InputIterator<Iter1>::value_type,
           InputIterator<Iter2>::value_type> }
  bool
  equal(Iter1 first1, Iter1 last1, Iter2 first2);
```

- `InputIterator` concept has **associated types**

- Supersede the use of traits
- Note: no `typename`!



# EqualityComparable vs. EqualityComparable2

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- ❑ `vector ==` requires `EqualityComparable`
- ❑ `std::equal` requires `EqualityComparable2`
- ❑ How are they related?



# EqualityComparable vs. EqualityComparable2

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- `vector ==` requires `EqualityComparable`
- `std::equal` requires `EqualityComparable2`
- How are they related?
  - `EqualityComparable2` is more general
  - We call `EqualityComparable` a **refinement** of `EqualityComparable2`.



# Refinement

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- A concept B **refines** concept A if B includes all of the requirements of A.
  - Akin to inheritance of abstract classes
  - `RandomAccessIterator` **refines** `BidirectionalIterator`
- A better way to define `EqualityComparable`:

```
template<typename T>
concept EqualityComparable : EqualityComparable2<T, T>
{ };
```





# InputIterator Concept

---

```
template<typeid Iter>
concept InputIterator : EqualityComparable<Iter>,
                      Assignable<Iter>,
                      CopyConstructible<Iter>
{
    typename value_type;
    typename difference_type;
    require Integral<difference_type>;
    const value_type& operator*(const Iter&);
    Iter& operator++(Iter&);
};
```



# “Make the hard things possible”

---

```
template<typeid X>
concept Sequence : Container<X>
{
    typename value_type;
    typename iterator;
    require ForwardIterator<iterator>,
        ForwardIterator<iterator>::value_type == value_type;

    template<typeid Iter>
    where {
        Convertible<InputIterator<Iter>::value_type,
            value_type> }
    X::X(Iter first, Iter last);
};
```



# Summary of features

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- Concepts
  - Refinement
  - Associated types
  - Pseudo-signatures
- Explicit models of concepts
- Where clauses describe requirements
  - Concept
  - Same-type
- Concept-based function selection



# Existing practice and the Standard Library

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# Standard Library $\Leftrightarrow$ Concepts

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- Requirements tables  $\Leftrightarrow$  □ Concepts
- Template type parameter names  $\Leftrightarrow$  □ Where clauses
- Loose syntactic requirements  $\Leftrightarrow$  □ Concepts + where clauses
- Traits  $\Leftrightarrow$  □ Associated types, model declarations



# Default Constructible

Utilities

Category: utilities

Concept

Component type: concept

## Existing practice

- Excerpt from SGI STL documentation:
- As written with our proposal:

```
template<typeid X>
concept DefaultConstructible
{
    X::X();
};
```

```
model DefaultConstructible<int> {};
model DefaultConstructible<vector<double> >
{ };
```

```
template<typename T>
    model DefaultConstructible<vector<T> >
{ };
```

### Description

A type is DefaultConstructible if it has a default constructor, that is, if it is possible to construct an object of that type without initializing the object to any particular value.

### Refinement of

### Associated types

### Notation

- x A type that is a model of DefaultConstructible
- x An object of type x

### Definitions

### Valid expressions

Name	Expression	Type requirements	Return type
Default constructor	x()		x
Default constructor	x x; <a href="#">[1]</a>		

### Expression semantics

Name	Expression	Precondition	Semantics	Postcondition
Default constructor	x()			
Default constructor	x x;			

### Complexity guarantees

### Models

- int
- [vector](#)<double>



# Requirements tables

[20.1.3/1] In the following Table 30,  $T$  is a type to be supplied by a C++ program instantiating a template,  $t$  is a value of type  $T$ , and  $u$  is a value of type  $\text{const } T$ .

```

template<typeid T>
concept CopyConstructible
{
    T::T(T&);
    T::T(const T&);
    T::~~T();
    T* operator&(T&);

    const T*
    operator&(const T&);
};
    
```

Table 30: CopyConstructible requirements

expression	return type	requirement
$T(t)$		$t$ is equivalent to $T(t)$
$T(u)$		$u$ is equivalent to $T(u)$
$t.\sim T()$		
$\&t$	$T^*$	denotes the address of $t$
$\&u$	$\text{const } T^*$	denotes the address of $u$



# Impact on the C++ community

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# Impact on users

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- Vast majority of code still works, unchanged
- Generic Programming becomes easier
  - Stronger typing for generic functions
  - Clean, concise error messages
  - Replace learning template tricks with learning Generic Programming
- Application code will need explicit model definitions
  - Not very often: the Standard Library handles most of them
  - Similar effort to inheriting abstract base classes



# Impact on Standard Libraries

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- Add `where` clauses as specified by the standard
- Option: conditionally enable `where` clauses for a single C++03/C++0x library
  - Tag dispatching & associated types a little harder
- Option: leave template parameters as `typename`
  - Checks uses of templates
  - ... but not definitions!
  - Allows clever optimizations in library implementations



# Impact on compilers

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- Several parts to implementation
  - Concepts, refinement, pseudo-signatures similar to class templates
  - Models similar to class template specializations
  - Type-checking is still two-phase lookup
  - Where clauses are quite simple
- We're addressing this on several fronts
  - Nontrivial portion of STL implemented in  $\mathcal{C}$
  - Prototype in GCC



# Summary of the proposal

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- We propose complete support for Generic Programming in C++ that:
  - Makes Generic Programming accessible
  - Reflects existing practice
  - Is expressive enough for the **entire** C++ Standard Library
  - Is implementable in current C++ compilers



# Explicit vs. Implicit Model Declarations

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# Definitions Recap

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- A **model declaration** states that a type or set of types meet the syntactic and semantic requirements of a concept.
- With **implicit** model declarations, the compiler performs **structural matches** to determine if a set of types model a concept.
- With **explicit** model declarations, the **user states** that a set of types model a concept.



# Implicit Model Declarations

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## Benefits

- User need not fully understand concepts
- Backward compatibility
- Matches what we do now in C++ (sometimes [\*])

## Problems

- Accidental conformance [\*]
- Implementation complexity [\*]



[\*] Indicates that we have examples for these points.

# Accidental Conformance

---

- Occurs when the semantics assumed due to a structural match are incorrect.
- Consider `istream_iterator`:
  - Structurally matches `ForwardIterator`.
  - Semantically matches `InputIterator`.
  - With implicit models, compiler can't catch this error.
- ```
vector<int> v(istream_iterator<int>(cin),  
            istream_iterator<int>());
```





# Accidental Conformance

---

- Occurs when the semantics assumed due to a structural match are incorrect.
- Consider `istream_iterator`:
  - Structurally matches `ForwardIterator`.
  - Semantically matches `InputIterator`.
  - With implicit models, compiler can't catch this error.
- ```
vector<int> v(istream_iterator<int>(cin),  
            istream_iterator<int>());
```
- Why doesn't this problem happen now?



# Accidental Conformance

---

- Occurs when the semantics assumed due to a structural match are incorrect.
- Consider `istream_iterator`:
  - Structurally matches `ForwardIterator`.
  - Semantically matches `InputIterator`.
  - With implicit models, compiler can't catch this error.
- ```
vector<int> v(istream_iterator<int>(cin),  
            istream_iterator<int>());
```
- Why doesn't this problem happen now?
  - `iterator_category` is an explicit model declaration!





# Nominal conformance in N1782

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- The refinement hierarchy uses nominal conformance
- Overloading is not structural (see Section 6.3 of N1782)



# Implementation Complexity

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- Structural matching requires SFINAE-like behavior
  - For arbitrary types & expressions
  - Compiler must quietly “back out” if structural match fails
  - Compiler initiates search for structural matches
- Implementors have claimed this is difficult



# Explicit Model Declarations

## Benefits

- Strong semantic guarantees
  - No accidental conformance
  - Better optimization [\*]
- Matches what we do now in C++ (sometimes)

## Problems

- Not fully backward-compatible [\*]
- Users must understand concepts [\*]
- Users must write model declarations



[\*] Indicates that we have examples for these points.

# Backward Compatibility

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- User will need to add model declarations
- Some ways to avoid these:
  - Where model declarations already exist as traits, we can create model templates (e.g., in the standard library)
  - Falling back to (weak) structural matching can be accomplished with model templates and metaprogramming.
  - Some concepts (e.g., DefaultConstructible) may be so basic that the compiler should add the models.



# “Porting” iterator\_traits

```
template<typename OldIter>
  where { Convertible<typename std::iterator_traits<OldIter>::category,
                  std::input_iterator_tag> }
  model InputIterator<OldIter> {};
```

```
template<typename OldIter>
  where { Convertible<typename std::iterator_traits<OldIter>::category,
                  std::output_iterator_tag> }
  model OutputIterator<OldIter> {};
```

```
template<typename OldIter>
  where { Convertible<typename std::iterator_traits<OldIter>::category,
                  std::forward_iterator_tag> }
  model ForwardIterator<OldIter> {};
```

```
template<typename OldIter>
  where { Convertible<typename std::iterator_traits<OldIter>::category,
                  Convertible<std::forward_iterator_tag>,
                  is_same<typename std::iterator_traits<OldIter>::reference,
                  typename std::iterator_traits<OldIter>::value_type&>::value }
  model MutableForwardIterator<OldIter> {};
```



# Better Optimization

---

- Explicit models are semantic guarantees
  - Needed by generic functions
  - Usable by compilers & optimizers
- Examples:
  - CopyConstructible/Assignable copy propagation
  - Parallelize RandomAccessIterator loops
  - VectorSpace loop fusion
- Is this feasible?
  - It's still somewhat of a research topic, but...
  - Could find some important optimization opportunities





# Users must understand concepts

- What if the user forgets a model declaration?
  - Compiler provides message like “type Foo does not model the InputIterator concept.”
  - User needs to know what that means.
- Mitigating factor: the compiler can suggest model declarations.

```
sort_list.cpp:7: error: no matching function for call to 'sort(std::_List_iterator<int>,
    std::_List_iterator<int>)'
/.../bits/stl_algo.h:2559: note: candidates are: void std::sort(_RandomAccessIterator,
    _RandomAccessIterator) [with _RandomAccessIterator = std::_List_iterator<int>]
    <failed requirements>
sort_list.cpp:7: note:   unable to locate a model
    'std::MutableRandomAccessIterator<std::_List_iterator<int> >'
sort_list.cpp:7: note:   for concept requirement
    'std::MutableRandomAccessIterator<_RandomAccessIterator>' (you may need to write a model
    definition)
```



# Our View

---

- Concepts have semantic requirements
  - We need users to state that their types meet these requirements
  - It's common practice to do so (e.g., iterators)
  - Optimization, simpler implementation just side benefits
- Backward compatibility is one-time hit
  - Big benefits once the jump is made
  - Compilers, libraries, and tools can help bridge the gap.



# Pseudo-signatures vs. Usage Patterns

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# Definitions recap

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- **Pseudo-signatures** look like function declarations or definitions, but match a class of functions.
- **Usage patterns** illustrate the syntax that generic functions will use, but match a class of functions.



# Semantics are the same

---

- One-sided leniency for both syntax kinds
  - Strict checking of concept uses in generic functions
  - “Convertible-to” okay in models (implicit and explicit)
  - Misunderstandings and prior issues obscure this
- Can both express concepts in the Standard Library
  - -> operator gives usage patterns some trouble
  - `OutputIterator` is hard on both syntax kinds



# Pseudo-signatures

---

- Declarations look like function declarations:

```
bool operator==(const T&, const U&);
```

```
bool operator!=(const T&, const U&);
```

```
template<typeid Iter>
```

```
    where {
```

```
        Convertible<InputIterator<Iter>::value_type,  
                    value_type> }
```

```
    X::X(Iter first, Iter last);
```



# Usage patterns

---

- Declare variables to be used in expressions:

```
T t;  
U u;  
Input_Iterator Iter;  
Iter i, j;  
static_assert  
    Convertible<Iter::value_type, value_type> {};
```

- Write expressions describing how objects can be used:

```
(bool) (t == u);  
(bool) (t != u);  
X(i, j);
```



# operator->

---

- Not possible in general for usage patterns (see N1782)
- Pseudo-signatures use, e.g.,:  
`T* operator->(const X&);`
- Impact: Probably need a built-in `Arrow` concept to support usage patterns





# Compound expressions

- Usage patterns can represent compound expressions:

```
T a, b, c;  
(T) (a * b + c);
```

- Internal type of `a * b` is unknown/irrelevant

- Pseudo-signatures require naming the type:

```
typename mul_type;  
mul_type operator*(T, T);  
T operator+(mul_type, T);
```

- `mul_type` can use a `decltype` default to save the user some effort



# Compound expressions: Impact

---

- Input and output iterator postincrement is specified as a compound expression
  - Just list the expression with usage patterns
  - Requires introduction of a hidden type for each concept
    - See N1758 for details



# Should models look like concepts?

- If a model looks like one of these:

```
static_assert FG<X> {  
    void f(X& a) { a.f(); }  
    void X::g() { g(*this); }  
};
```

```
model FG<X> {  
    void f(X& a) { a.f(); }  
    void X::g() { g(*this); }  
};
```

- What should the concept look like?

```
concept FG<class T> {  
    T a;  
    f(a);  
    a.g();  
};
```

```
template<typeid T>  
concept FG {  
    void f(X& a);  
    void X::g();  
};
```



# Expression templates

---

- Expression templates require lots of “hidden” types
  - One for each subexpression
- Impact:
  - Usage patterns: list every expression in the function body as a usage pattern
  - Pseudo-signatures: add pseudo-types for each subexpression
- Distinction between “opaque” and “non-opaque” types in N1758 offers one solution.



# Syntax: Pseudo-signatures

---

- Pros:
  - Match free & member function syntax
  - Model syntax matches concept syntax
  - Concepts resemble abstract classes
  - Very little new parser technology
- Cons:
  - More verbose than usage patterns
  - Compound expressions are more painful
  - Standard Library (and other generic libraries) use usage patterns/valid expressions



# Syntax: Usage patterns

---

## □ Pros:

- Concise
- Similar to existing requirements tables
- Very little new parser technology

## □ Cons:

- Usage patterns look different from the things they describe (no cut 'n' paste coding)
- Need built-in `Arrow` concept
- Determining which declarations declare values vs. types can be confusing.



# Our View

---

- It's just a syntax issue
- Which is “more readable”?
  - It's a toss up: good and bad examples for both
  - We like how pseudo-signatures look like the function declarations they match
    - How does one implement a new model of a concept? Copy 'n' paste!
    - No “hidden” template requirements



# Associated Types

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# Where do associated types live?

---

- N1758: As associated types of concepts
  - `ForwardIterator<Iter>::value_type`
- N1782: As member types of types involved in concepts
  - `Iter::value_type`



# Ambiguity with member types

```
concept Callable1<F,T1> {
    typename F::result_type;
    F f; T1 t1;
    (F::result_type) (f(t1));
};

struct negate {
    template<typename T> operator() (T x) { return -x; }
};

static_assert Callable1<negate, int>
    { typedef int negate::result_type; };

static_assert Callable1<negate, float>
    { typedef float negate::result_type; };

// (what is negate::result_type?)
```



# Our View

---

- Using member types can be awkward
  - Potential ambiguities with syntax
  - Odd to add member types to non-classes
  - There isn't always a “main type” to hang on
- Associated types are naturally part of concepts
  - Multiple types of a concept participate to produce associated types
  - Traits use this same level of indirection



# Semi-structured additional slides

---

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# std::distance with concepts

```
template<typeid Iter> where { InputIterator<Iter> }  
InputIterator<Iter>::difference_type  
distance(Iter first, Iter last)  
{  
    InputIterator<T>::difference_type result = 0;  
    for (; first != last; ++first)  
        ++result;  
    return result;  
}
```

## □ Things to notice:

- We've added more to the declaration
- The body really hasn't changed



# std::distance with concepts

```
template<typeid Iter> where { InputIterator<Iter> }
InputIterator<Iter>::difference_type
distance(Iter first, Iter last)
{
    InputIterator<T>::difference_type result = 0;
    for (; first != last; ++first)
        ++result;
    return result;
}
template<typeid Iter>
    where { RandomAccessIterator<Iter> }
RandomAccessIterator<Iter>::difference_type
distance(Iter first, Iter last)
{
    return last - first;
}
```



# Concepts

```
template<typename Iter> where { InputIterator<Iter> }  
InputIterator<Iter>::difference_type  
distance(Iter first, Iter last)  
{  
    InputIterator<T>::difference_type result = 0;  
    for (; first != last; ++first)  
        ++result;  
    return result;  
}
```

- **A concept is a set of requirements**
  - Syntactic: functions, operators, types
  - Semantic: what functions do, function complexity
  - Like requirements tables in the standard



# InputIterator Concept

```
template<typeid Iter>
concept InputIterator : EqualityComparable<Iter>,
                      Assignable<Iter>,
                      CopyConstructible<Iter>
{
    typename value_type;
    typename difference_type;

    const value_type& operator*(const Iter&);
    Iter& operator++(Iter&);
};
```

## □ Things to note:

- Has types `value_type` and `difference_type`
- Has pseudo-signatures for operators `++` and `*`
- **Refines** `EqualityComparable`, `Assignable`, and `CopyConstructible`





# Opaque template parameters

```
template<typeid Iter> where { InputIterator<Iter> }  
InputIterator<Iter>::difference_type  
distance(Iter first, Iter last)  
{  
    InputIterator<T>::difference_type result = 0;  
    for (; first != last; ++first)  
        ++result;  
    return result;  
}
```

- ❑ typename eliminates type checking: you can do anything to typename types.
- ❑ typeid provides type checking: you can do nothing but what you require
- ❑ Type checking has two phases already!



# where clauses

```
template<typename Iter> where { InputIterator<Iter> }  
InputIterator<Iter>::difference_type  
distance(Iter first, Iter last)  
{  
    InputIterator<Iter>::difference_type result = 0;  
    for (; first != last; ++first)  
        ++result;  
    return result;  
}
```

- where clauses constrain templates
  - Example: Iter must **model** the InputIterator concept
  - User must pass in types that meet these requirements
  - Implementor must only use types and operations mandated by the requirements



# Associated types

```
template<typename Iter> where { InputIterator<Iter> }  
InputIterator<Iter>::difference_type  
distance(Iter first, Iter last)  
{  
    InputIterator<T>::difference_type result = 0;  
    for (; first != last; ++first)  
        ++result;  
    return result;  
}
```

- Associated types are additional types used to specify a concept
  - Remember `difference_type` from `InputIterator`?
  - Supersedes traits
  - Note: no `typename`!



# How do we call distance?

---

- Just like we always have:

```
list<int> l; // initialize l
cout << "Length = "
      << distance(l.begin(), l.end()) << endl;
```

- Most uses of concepts will be invisible
  - They provide better type safety
  - It's easier to write correct templates
  - Standard library provides its own models
  - Compiler will provide some concepts “for free.”



# More same-type constraints

## □ Consider `std::merge`:

```
template<typename Iter1, typename Iter2, typename OIter,
         typename T>
where {
    InputIterator<Iter1>::value_type == T,
    InputIterator<Iter2>::value_type == T,
    Convertible<OutputIterator<OIter>::value_type, T>,
    StrictWeakOrdering<T> }
OIter merge(Iter1 first1, Iter1 last1,
            Iter2 first2, Iter2 last2, OIter out);
```



# Sequence requirements

Table 68: Sequence requirements (in addition to container)

| expression           | return type | assertion/note<br>pre/post-condition                                                        |
|----------------------|-------------|---------------------------------------------------------------------------------------------|
| X(n, t)<br>X a(n, t) |             | post: size() == n<br>constructs a sequence with n copies of t                               |
| X(i, j)<br>X a(i, j) |             | post: size() == distance between i and j<br>constructs a sequence equal to the range [i, j) |
| a.insert(p,t)        | iterator    | inserts a copy of t before p                                                                |
| a.insert (p,n,t)     | void        | inserts n copies of t before p                                                              |
| a.insert (p,i,j)     | void        | pre: i and j are not iterators into a.<br>inserts copies of elements in [i, j) before p     |
| a.erase(q)           | iterator    | erases the element pointed to by q                                                          |
| a.erase(q1,q2)       | iterator    | erases the elements in the range [q1, q2).                                                  |
| a.clear()            | void        | erase(begin(), end())<br>post: size() == 0                                                  |
| a.assign(i,j)        | void        | pre: i, j are not iterators into a.<br>Replaces elements in a with a copy of [i, j).        |
| a.assign(n,t)        | void        | pre: t is not a reference into a.<br>Replaces elements in a with n copies of t.             |



# Why propose language support?

---

- Generic Programming is gaining acceptance
  - C++ Standard Library is prototypical example
  - Still too much work to use GP
  - Java, C# now support templates/generics
- Conventions are okay, but language support is better
  - The ideas of GP can be expressed more clearly
  - Better tool support
- We are at the *tipping point*



# Revisiting vector ==

- How does this type-check?:

```
template<typeid T, typeid Alloc>
  where { EqualityComparable<T> }
  bool operator==(const vector<T, Alloc>& x,
                  const vector<T, Alloc>& y)
  {
    return x.size() == y.size()
           && equal(x.begin(), x.end(), y.begin());
  }
```

- The only operations on T objects are in `std::equal`
- `std::equal` requires `EqualityComparable2`
  - Okay, since `EqualityComparable` implies `EqualityComparable2`
- `std::equal` requires `InputIterators`
  - Okay, since `vector<T>::const_iterator` models `RandomAccessIterator`.





# “Small” example

---

```
template<typeid T, int N>
concept Small
{
    require sizeof(T) <= N;
};
```

```
template<typename T, int N>
    where {sizeof(T) <= N }
    model Small<T, N> {};
```

```
template<typename T> where {Small<T,200> } void f(T&);
template<typename T> void f(T&);
```

```
template<typename T>
    void foo(const T& t)
    {
        f(t);
    }
```

