`std::generator`: Synchronous Coroutine Generator for Ranges

Abstract

We propose a standard library type `std::generator` which implements a coroutine generator that models `std::ranges::input_range`.

Acknowledgements

I'd like to thank Lewis Baker and Corentin Jabot, whose P2168 [3] I looted shamelessly for this proposal. This paper is presented as a new revision of their unpublished D2168R4 for continuity of design.

Revisions

P2502R1

- The type `generator::yielded` is now publicly-visible and not exposition-only per 2021-12-13 LEWG discussion.

- Feature-test macro insertion is now ordered properly with the other wording.

- Coroutines for `generator<T&&>` can now yield lvalues; the generated element is an xvalue denoting a copy of the lvalue. This saves folks the trouble of writing `co_yield auto(lvalue)` instead of `co_yield lvalue`.

- Reorder `generator`'s template parameters (again, for the last last time). Coroutine Task Force discussions convinced all participants (including the author) that it was too big a break with the rest of the Standard Library to have the `Allocator` not be the final template parameter.

- Simplify `elements_of` into a vanilla two-element `struct`.

- Update implementations accordingly.

- Per 2022-01-025 LEWG direction, `generator<T>` now behaves like `generator<T&&>` instead of like `generator<const T&>` per the design in P2529R0 [9].

P2502R0

- Reorder `generator`'s template parameters. This allows the reference type to be more easily defaulted to a true reference, while still respecting requests for differing value and reference types. This preserves the previous design's ease-of-use, while providing full generality.

- Remove concerns about the $\mathcal{O}(1)$ destruction requirement for `view`, which has been relaxed by P2415R2 "What is a view?" [7].

D2168R4

- Wording improvements

P2168R3

- Wording improvements

P2168R2

- Some wording fixes
- Improve the section on allocator support
- Updated implementation

P2168R1

- Add benchmarks results and discussion about performance
- Introduce `elements_of` to avoid ambiguities when a generator is convertible to the reference type of the parent generator.
- Add allocator support
- Symmetric transfer works with generators of different value / allocator types
- Remove *iterator*::`operator->`
- Put `generator` in a new `<generator>` header.
- Add an other example to motivate the `Value` template parameter

Example

```
std::generator<int> fib() {
    auto a = 0, b = 1;
    while (true) {
        co_yield std::exchange(a, std::exchange(b, a + b));
    }
```

```
    }

    int answer_to_the_universe() {
        auto rng = fib() | std::views::drop(6) | std::views::take(3);
        return std::ranges::fold_left(std::move(range), 0, std::plus{});
    }
```

Motivation

C++ 20 had very minimalist library support for coroutines. Synchronous generators are an important use case for coroutines, one that cannot be supported without the machinery presented in this paper. Writing an efficient recursive generator is non-trivial; the standard should provide one.

Design

While the proposed `std::generator` interface is fairly straightforward, a few decisions are worth pointing out.

`input_view`

`std::generator` is a move-only `view` which models `input_range` and has move-only iterators. This is because the coroutine state is a unique resource (even if the coroutine *handle* is copyable).

Header

Multiple options are available as to where to put the `generator` class.

- `<coroutine>`, but `<coroutine>` is a low level header, and `generator` depends on bits of `<type_traits>` and `<iterator>`.

- `<ranges>`

- A new `<generator>`

This paper uses a new `<generator>` header since P2168R3 did so, and LEWG has provided no guidance to do otherwise. We do note on our MSVC STL branch implementation that `#include<ranges>` includes 52.6k lines of code, and `#include<generator>` 53.3k lines. (Note that `<generator>` is specified to include both `<ranges>` and `<coroutines>`.) Defining `generator` in `<ranges>` together with a `#include<coroutine>` would penalize people who want `<ranges>` but not `generator` by about 610 LoC.

Reference type

`generator` has 3 template parameters: `generator<R, V = void, Allocator = void>`

From `R` and `V`, we derive types:

```
using Value = conditional_t<is_void_v<V>, remove_cvref_t<R>, V>;}
using Reference = conditional_t<is_void_v<V>, R&&, R>;
using Yielded = conditional_t<is_reference_v<Reference>, Reference, const Reference&>;
```

- `Value` is a cv-unqualified object type that specifies the value type of the `generator`'s range and iterators,

- `Reference` specifies the reference type (not necessarily a core language reference type) of the `generator`'s range and iterators, and

- `Allocator` is the type of allocator used for the coroutine state, which can be `void` to type-erase any allocator specified as a coroutine argument, defaulting to `allocator<byte>` when none is specified.

- `Yielded` (necessarily a reference type) is the type of the parameter to the primary overload of `yield_value` in the generator's associated promise type.

generator<meow>

Our expectation is that 98% of use cases will need to specify only one parameter. The resulting generator:

- has a value type of `remove_cvref_t<meow>`

- has a reference type of `meow`, if it is a reference type, or `meow&&` otherwise,

- expects `co_yield` to appear in the body of the generator with operands that are convertible to the reference type, and

- can use any allocator (via type-erasure) defaulting to `allocator<byte>`.

This avoids the performance pitfall from earlier revisions of the proposal that used the first argument type directly as reference type; users who naively chose `generator<std::string>` got an iterator that produces independent copies of the yielded value on every dereference, when they may have been satisfied by yielding a reference to the same constant value.

As a nod to ease of use, and to more readily reflect the semantics of prvalue-returning functions, we allow yielding lvalues from generators with an rvalue reference type when the lvalue is copyable. The result is that the yielded lvalue is copied into temporary storage, and an iterator for the generator will return an xvalue denoting that copy when dereferenced.

generator<meow, woof>

For the rare user who needs `generator` to step outside the box and use a proxy reference type, or who needs to generate a range whose iterators yield prvalues for whatever reason, we have two-argument `generator`. If `woof` is `void`, this is `generator<meow>`. Otherwise, the resulting generator:

- has a value type of `woof`,

- has a reference type of `meow`,

- expects `co_yield` to appear in the body of the generator with operands that are convertible to `meow`, if it is reference type, and otherwise `const meow&`, and

- can use any allocator (via type-erasure) if `woof` is `void`, or otherwise can use any allocator convertible to `woof`.

Your iterators can yield a prvalue, but it must be a prvalue of `copy_constructible` type so a copy of the operand of a single `co_yield` can be returned multiple times from repeated dereferences of the same iterator value.

generator<meow, woof, quack>

For use cases that want to specify an allocator type statically so they need not constantly pass pairs of `allocator_arg, my_allocator` arguments to every coroutine, we have three-argument `generator`. The resulting generator can use any allocator convertible to `quack`, defaulting to a default-constructed `woof` if it is `default_initializable`. The value, reference, and yielded types are as described for `generator<meow>` if `woof` is `void`, and as described for `generator<meow, woof>` otherwise.

Our expectation is that libraries that wish to declare many functions with the same statically-specified allocator type will define a template alias like

```
template <class R, class V = void>
using my_generator = generator<R, V, my_allocator>;
```

to ease declarations. In practice, this should mean `generator<R, void, my_allocator>` (with an explicit second argument of `void`) appears quite rarely in real code.

Allocator = void

In the 2022-01-13 Coroutine Task Force meeting, Lewis Baker brought up that LEWG feedback on some revision of P2168 had expressed dissatisfaction with the use of `void` to indicate "type-erase any allocator." We discussed some alternatives:

- Devise a new type for this specific purpose. We felt that single-use tag types have low return-on-investment; every such tag type consumes a few more bytes of memory in nearly every compilation, and a few more neurons of memory in every programmer.

- Reuse an existing type for this purpose. `void` has the advantage that it can never be confused for a conforming allocator type, even if it someday does become a regular type. `void` requires no predeclaration, and no memory is is expended to store its definition.

  `monostate` was another suggestion that was universally reviled; that type already has a very clear purpose in the Standard Library that we didn't feel inclined to muddle.

  No one suggested another type that we found to be an improvement.

Our conclusion was that we didn't find the use of `void` for this purpose confusing, we were happy to keep it given that more-recent LEWG discussion had resulted in no further feedback on this topic.

Obsolete discussion about reference specification

[*Note:* Before P2502R0, `generator`'s first parameter `Type` denoted the reference type of the range / iterators, and the value type was defaulted to `remove_cvref_t<Type>`. The following sections of design discussion from that era are preserved here. — *end note*]

In earlier versions of this paper, the reference type was exactly the first template parameter. This had the advantage of being simple. But it was a terrible performance trap:

Consider the behavior of the following code assuming the reference type is exactly the first template argument:

```
std::generator<std::string> f() {
    std::string hello = "hello";
    co_yield hello;   // 0 or 1 copy depending on implementation
    co_yield "Hello"; // 1 copy (conversion from const char* to std::string)
}

for (auto&& str : f()) {} // 1 copy (*it returns std::string)
```

Of course the solution, which we advocated for, is for the user to manually specify an explicit reference type:

```
    generator<const std::string&> f() {
        std::string hello = "hello";
        co_yield hello;   // 0 or 1 copy depending on implementation
        co_yield "Hello"; // 1 copy (conversion from const char* to std::string)
    }

    for (auto&& str : f()) {} // 0 copy
```

This works, can be explained, and is even logical. You get what you asked for. It is nonetheless surprising for non-experts that using the simple `generator<string>` would create 2 copies per `co_yield`.

To hope users would not routinely forget to use a reference type when using `std::generator` calls for a heaping barrel of optimism.

We later proposed that for a `generator<T>`, its reference type be `conditional_t<is_reference_-v<T>, T, const T&>`.

First parameter	reference type	default value	can yield mutable lvalue ref?
int	const int&	int	No
const int&	const int&	int	No
int&	int&	int	Yes
int&&	int&&	int	No
const int&&	const int&&	int	No

6

Attempts have been made to characterize the exact relations between reference, value, storage, and `co_yield` exception types and categories. Ultimately, a simpler mental model is to characterize what expressions can be yielded for a given reference type and how many copies are made for each scenario.

First parameter	`co_yield const T&`	`co_yield T&`	`co_yield T&&`	`co_yield U&&`
`T`	0	0	0	1
`const T&`	0	0	0	1
`T&`	Ill-formed	0	Ill-formed	Ill-formed
`T&&`	Ill-formed	Ill-formed	0	1
`const T&&`	Ill-formed	Ill-formed	0	1

In this table, we see that only `co_yield` that requires conversion incurs copy, which is expected. Coroutines guarantee that the yielded expression exceeds the lifetime of the `co_yield` expression, so `generator` can usefully store a pointer to the object denoted by a yielded xvalue.

`co_yield` expressions involving conversion can store the yielded value in an awaiter. The type of the stored expression is the reference type with its reference qualifiers stripped, but that is an implementation detail that is not observable and is therefore of limited interest. Of course, that type needs to be constructible from yielded values.

Besides the `T` case, this behaves very much like returning from a function that is intended.

Move-only and immovable types

LEWG was interested in how this works with `generator` of move-only and immovable types.

First parameter	`co_yield const T&`	`co_yield T&`	`co_yield T&&`
`move_only`	0	0	0
`const move_only&`	0	0	0
`move_only&`	Ill-formed	0	Ill-formed
`move_only&&`	Ill-formed	Ill-formed	0
`const move_only&&`	Ill-formed	Ill-formed	0
`immovable`	0	0	0
`const immovable&`	0	0	0
`immovable&`	Ill-formed	0	Ill-formed
`immovable&&`	Ill-formed	Ill-formed	0
`const immovable&&`	Ill-formed	Ill-formed	0

As that table shows, these types work exactly like other types. However, to be able to move from a move only reference type, the coroutine has to explicitly state so:

```cpp
auto f = []() -> std::generator<move_only> { co_yield move_only{}; }();
for (auto&& x : f) {
    move_only mo = std::move(x); // ill-formed, decltype(x) is const move_only&
}

auto f = []() -> std::generator<move_only&&> { co_yield move_only{}; }();
for (auto&& x : f) {
```

```
    move_only mo = x; // ok
}

auto f = []() -> std::generator<move_only&> { move_only m; co_yield m; }();
for (auto&& x : f) {
    move_only mo = std::move(x); // dicey but okay
}
```

Potential downsides

```
auto f = []() -> std::generator<MyType> {
    MyType t;
    co_yield std::move(t);
}();
```

In the example above `std::move` doesn't move. Arguably more than usual. Indeed the code expands to something similar to:

```
auto&& __temp = std::move(t);
yield_value(_temp); // <=> promise.value = std::addressof(__temp); // no move
```

Of course, a move would not have occurred for a `std::generator<const MyType&>` either as these things are identical. It might be suprising? The only way to avoid that is to create temporary value for rvalue reference, which would force a move to actually occurs, at the cost of performance.

Alternatives considered

Mandating a reference as the first parameter   We could make `generator<int>` ill-formed and force people to specify a reference type like `generator<const int&>`. We do not think this is very user-friendly, given that we can provide a reasonable default.

We rejected this option.

Using `T&` as the default   There are two issues with mutable references:

- They are mutable (They allow mutating the coroutine frame), which would be an *interesting* default.
- They are very restrictive as to the set of `co_yield` expression allowed with them.

We rejected this option.

Using `T&&` as the default   This avoids a copy when doing `auto object = *it` (where it is a `std::generator::iterator`), but it is easy to misuse, consider:

```
auto f = [] -> std::generator<std::string> { co_yield "footgun"; }();
for (auto&& x : f | std::views::filter([](std::string s) { return s.size() > 0; })) {
    std::cout << x << '\n'; // outputs a single empty line
}
```

We rejected this option. [*Note:* Nevertheless, we came back to this option after much discussion of pros and cons in both LEWG and several additional meetings if a Coroutine Task Force formed by LEWG Chair specifically to build consensus. *— end note*]

Doing something clever for move-only types   We considered returning `T&` for move_only types so that they can be moved from by default. We realized this was too clever and inconsistent. Notably, adding a copy constructor to `T` would change the meaning of the code.

We rejected this option.

Doing something clever for reference types   By default `generator<reference_wrapper<T>>` could yield `reference_wrapper<T>` has that is already a "reference-like" type. However, no other view does that, "reference-like" is fuzzily defined, and this would probably cause more trouble than it's worth.

We rejected this option.

Keeping the D2168R4 design   Returning values has the potential to severely impact performance, is inconsistent with other views, and is not necessary. It also did not work with move-only types.

The change, along with an implementation strategy described in the "How to store the yielded value in the promise type?" guarantees that no copy needs to be made if the reference and yielded types are the same (with qualifiers stripped).

We think this new approach keeps the simplicity of the original design, improves performance, and works with more types.

Thank you LEWG, and in particular Mathias, for highlighting these concerns!

Separately specifyable Value Type

This proposal supports specifying both the "yielded" type, which is the iterator's reference type (not required to be a reference) and its corresponding value type. This allow ranges to handle proxy types and wrapped `reference`, like this implementation of `zip`:

```cpp
namespace ranges = std::ranges;

template<ranges::input_range Rng1, ranges::input_range Rng2>
std::generator<
    std::tuple<ranges::range_reference_t<Rng1>, ranges::range_reference_t<Rng2>>
    std::tuple<ranges::range_value_t<Rng1>, ranges::range_value_t<Rng2>>>
zip(Rng1 r1, Rng2 r2) {
    auto it1 = ranges::begin(r1);
    auto it2 = ranges::begin(r2);
    auto end1 = ranges::end(r1);
    auto end2 = ranges::end(r2);
    for (; it1 != end1 && it2 != end2; ++it1, ++it2) {
        co_yield {*it1, *it2};
```

```
        }
}
```

In this second example, using `string` as value type ensures that calling code can take the necessary steps to make sure iterating over a generator would not invalidate any of the yielded values.

```cpp
// Yielding string literals : always fine
std::generator<std::string_view, std::string_view> string_views() {
    co_yield "foo";
    co_yield "bar";
}

std::generator<std::string_view, std::string> strings() {
    co_yield "start";
    std::string s;
    for (auto sv : string_views()) {
        s = sv;
        s.push_back('!');
        co_yield s;
    }
    co_yield "end";
}

// conversion to a vector of strings
// If the value_type was string_view, it would convert to a vector of string_view,
// which would lead to undefined behavior operating on elements of v that were
// invalidated while iterating through the generator.
auto v = std::ranges::to<vector>(strings()); // (P1206R3 [4])
```

How to store the yielded value in the promise type?

There are multiple implementation strategies possible to store the value in the generator. An early revision of this paper always stored a copy of the yielded value, leading to an extra copy. Later revisions supported storing the yielded value in an awaitable object returned from the promise's `yield_value` function.

However, the object denoted by a glvalue yield expression is guaranteed to live until the coroutine resumes. We can take advantage of that fact by storing only a pointer to the denoted object in the promise, if the result of dereferencing that pointer is convertible to the generator's reference type. We guarantee this is the case by providing a `yield_value` whose parameter type is always a reference type (`conditional_t<is_reference_v<Reference>, Reference, const Reference&>`). This forces any conversions to happen inside the coroutine itself, yielding a temporary glvalue that can later be dereferenced to an lvalue which is trivially `static_cast`ed to `Reference` in the iterator's `operator*`.

A drawback of this solution is that the yielded value is only destroyed at the end of the full expression in which `co_yield` appears, so given

```cpp
(co_yield x, co_yield y); // x is destroyed after y is yielded.
```

We think this is a reasonable tradeoff given that this approach minimizes the number of copies must be made of the yielded value. We force the coroutine to materialize the element to be yielded, but after doing so can cleanly pass a reference to that element through the coroutine and iterator machinery and directly to consuming code.

Recursive generator

A "recursive generator" is a coroutine that supports the ability to directly `co_yield` a generator of the same type as a way of emitting the elements of that `generator` as elements of the current `generator`.

Example: A `generator` can `co_yield` other generators of the same type

```cpp
std::generator<const std::string&> delete_rows(std::string table, std::vector<int> ids) {
    for (int id : ids) {
        co_yield std::format("DELETE FROM {0} WHERE id = {1};", table, id);
    }
}

std::generator<const std::string&> all_queries() {
    co_yield std::ranges::elements_of(delete_rows("user", {4, 7, 9 10}));
    co_yield std::ranges::elements_of(delete_rows("order", {11, 19}));
}
```

Example: A `generator` can also be used recursively

```cpp
using namespace std;

struct Tree {
    Tree* left;
    Tree* right;
    int value;
};

generator<int> visit(Tree& tree) {
    if (tree.left) co_yield ranges::elements_of(visit(*tree.left));
    co_yield tree.value;
    if (tree.right) co_yield ranges::elements_of(visit(*tree.right));
}
```

In addition to being more concise, the ability to directly yield a nested generator has some performance benefits compared to iterating over the contents of the nested generator and manually yielding each of its elements.

Yielding a nested `generator` allows the consumer of the top-level coroutine to directly resume the current leaf generator when incrementing the iterator, whereas a solution that has each generator manually iterating over elements of the child generator requires O(depth) coroutine resumptions/suspensions per element of the sequence.

Example: Non-recursive form incurs O(depth) resumptions/suspensions per element and is more cumbersome to write:

```cpp
using namespace std;

generator<int> slow_visit(Tree& tree) {
    if (tree.left) {
        for (int x : ranges::elements_of(visit(*tree.left)))
            co_yield x;
    }
    co_yield tree.value;
    if (tree.right) {
        for (int x : ranges::elements_of(visit(*tree.right)))
            co_yield x;
    }
}
```

Exceptions that propagate out of the body of nested `generator` coroutines are rethrown into the parent coroutine from the `co_yield` expression rather than propagating out of the top-level `iterator::operator++()`. This follows the mental model that `co_yield someGenerator` is semantically equivalent to manually iterating over the elements and yielding each element.

For example: `nested_ints()` is semantically equivalent to `manual_ints()`

```cpp
std::generator<int> might_throw() {
    co_yield 0;
    throw some_error{};
}

std::generator<int> nested_ints() {
    try {
        co_yield std::ranges::elements_of(might_throw());
    } catch (const some_error&) {}
    co_yield 1;
}

// nested_ints() is semantically equivalent to the following:
std::generator<int> manual_ints() {
    try {
        for (int x : might_throw()) {
            co_yield x;
        }
    } catch (const some_error&) {}
    co_yield 1;
}

void consumer() {
    for (int x : nested_ints()) {
        std::cout << x << " "; // outputs 0 1
    }

    for (int x : manual_ints()) {
        std::cout << x << " "; // also outputs 0 1
    }
}
```

`std::ranges::elements_of`

`ranges::elements_of` is a utility function that prevents ambiguity when a nested generator type is convertible to the value type of the present generator

```cpp
generator<int> f()
{
    co_yield 42;
}

generator<any> g()
{
    co_yield f(); // should we yield 42 or generator<int> ?
}
```

To avoid this issue, we propose that:

- `co_yield <expression>` yields the value directly, and

- `co_yield elements_of(<expression>)` yields successive elements from the nested generator.

For convenience, we further propose that `co_yield elements_of(x)` be extended to support yielding the values of arbitrary ranges beyond generators, ie

```cpp
std::generator<int> f()
{
    std::vector<int> v = /*... */;
    co_yield std::ranges::elements_of(v);
}
```

Symmetric transfer

The recursive form can be implemented efficiently with symmetric transfer. Earlier works in [CppCoro] implemented this feature in a distinct `recursive_generator` type.

However, it appears that a single type is reasonably efficient thanks to HALO optimizations and symmetric transfer. The memory cost of that feature is two extra pointers per generator[1]. It is difficult to evaluate the runtime cost of our design given the current coroutine support in compilers. However our tests show no noticeable difference between a `generator` and a `recursive_generator` which is called non-recursively. It is worth noting that the proposed design makes sure that HALO [8] optimizations are possible.

While we think a single `generator` type is sufficient and offers a better API, there are three options:

- A single `generator` type supporting recursive calls (this proposal).

---

[1]The two pointers in our implementation have non-overlapping active times; we believe the pair can be optimized into a single pointer's space with some bit hacking to store a discriminator in the unused lower bits.

- A separate type `recursive_generator` that can yield values from either a `recursive_generator` or a `generator`. That may offer very negligible performance benefits, same memory usage.

- A separate recursive_generator type which can only yield values from other `recursive_generator`s.

  That third option would make the following ill-formed:

  ```
  generator<int> f();
  recursive_generator<int> g() {
      co_yield f(); // incompatible types
  }
  ```

  Instead you would need to write:

  ```
  recursive_generator<int> g() {
      for (int x : f()) co_yield x;
  }
  ```

  Such a limitation can make it difficult to decide at the time of writing a generator coroutine whether or not you should return a `generator` or `recursive_generator` as you may not know at the time whether or not this particular generator will be used within `recursive_generator` or not.

  If you choose the `generator` return-type and then later someone wants to yield its elements from a `recursive_generator` then you either need to manually yield its elements one-by-one or use a helper function that adapts the `generator` into a `recursive_generator`. Both of these options can add runtime cost compared to the case where the generator was originally written to return a `recursive_generator`, as it requires two coroutine resumptions per element instead of a single coroutine resumption.

  Because of these limitations, we are not recommending this approach.

Symmetric transfer is possible for different generator types as long as the `reference` type is the same, aka, different value type or allocator type does not preclude symmetric transfer (see the section on allocators).

Allocator support

In line with the design exploration done in section 2 of P1681R0 [6], `std::generator` supports both stateless and stateful allocators and strives to minimize the interface verbosity for stateless allocators by templating both the generator itself and the `promise_type`'s `new` operator on the allocator type. Details for this interface are found in P1681R0 [6].

`coroutine_parameter_preview_t` such as discussed in section 3 of P1681R0 [6] has not been explored in this paper.

```cpp
std::generator<int> stateless_example() {
    co_yield 42;
}

template <class Allocator>
std::generator<int> allocator_example(std::allocator_arg_t, Allocator alloc) {
    co_yield 42;
}

my_allocator<std::byte> alloc;
input_range auto rng = allocator_example(std::allocator_arg, alloc);
```

The proposed interface requires that, if an allocator is provided, it is the second argument to the coroutine function, immediately preceded by an instance of `std::allocator_arg_t`. This approach is necessary to distinguish the allocator desired to allocate the coroutine state from allocators whose purpose is to be used in the body of the coroutine function. The required argument order might be a limitation if any other argument is required to be the first. However, we cannot think of any scenario where that would be the case.

We think it is important that all standard and user coroutine types can accommodate similar interfaces for allocator support. In fact, the implementation for that allocator support can be shared amongst `generator`, `lazy`, and other standard types.

By default `std::generator` type erases the allocator type, and uses `std::allocator` unless an allocator is provided to the coroutine function. Then:

Type erased allocator(default)

```cpp
template <class Allocator>
std::generator<int> f(std::allocator_arg_t, Allocator alloc) {}

f(std::allocator_arg, my_alloc{});
```

Returns a generator of type `std::generator<int, void, void>` where the final `void` denotes that the allocator is type erased. The allocator is stored in the same allocation as the coroutine state if it is stateful or not default constructible; a pointer is always stored so that the `deallocate` method of the type erased allocator can be called.

No allocator

```cpp
std::generator<int> f() {}
f();
```

Again, returns a generator of type `std::generator<int, void, void>` where the final `void` denotes that the allocator is type erased. A pointer is stored so that the `deallocate` method of the type-erased allocator can be called, but the default allocator (`std::allocator`) need not be stored since it is stateless.

Explicit stateless allocator

```cpp
std::generator<int, void, std::stateless_allocator<int>> f() {}
f();
```

No extra storage is used for the allocator because it is stateless.

Explicit stateful allocator

```
std::generator<int, void, some_stateful_allocator<int>>
    f(std::allocator_arg_t, some_stateful_allocator<int> alloc) {}
f(std::allocator_arg, some_allocator); // must be convertible to some_stateful_allocator
```

The allocator is copied in the coroutine state.

Can we postpone adding allocator support?

A case can be made that allocator support could be added to `std::generator` later. However, because the proposed design has the allocator as a template parameter, adding allocator after `std::generator` ships would represent an ABI break. We recommend that we add allocator support as proposed in this paper now and make sure that the design remains consistent as work on `std::lazy` is made in this cycle. However, it would be possible to extend support for different mechanisms (such as presented in section 3 of P1681R0 [6] later.

Interaction of symmetric transfer and allocator support

The allocator must necessarily be part of a coroutine's promise type since implementations query the promise for allocation functions. Nonetheless, it would seem silly for a generator to be unable to nest another generator with identical element type but differing allocator. For that matter, even differing value types shouldn't be problematic: the only interface between the generator and the coroutine it wraps that differs depending on the type arguments to `generator` is `yield_value`. Ideally, generators would be able to recurse into other generators whose `yield_value` has the same parameter type even if all three template arguments to `generator` differ.

Our implementation uses a base class to implement the non-allocation behaviors for `generator`'s promise type so that generators with different allocator types can yield each other. Doing so, however, requires that we partially erase the type of a `coroutine_handle` so we can resume it later knowing only that its promise type derives from a particular base.

There are at least two ways to implement this partial type erasure:

- Storing a pointer in the common base to a component with full type knowledge, which can then resume the targeted coroutine,

- Relax the preconditions on some of the `coroutine_handle` functions to allow conversion from `coroutine_handle<void>` to `coroutine_handle<T>` when the source's corresponding `address()` value was obtained from a `coroutine_handle` referring to a coroutine whose promise object is pointer-interconvertible with an object of type `T`.

Our current plan is to standardize the intent to allow yielding nested generators with different allocator and value types, leaving the details of the implementation unspecified, and to later separately propose the changes to `coroutine_handle` that enable that implementation to be maximally efficient.

Implementation and experience

`generator` has been provided as part of cppcoro and folly. However, cppcoro offers a separate `recursive_generator` type, which is different than the proposed design.

Folly uses a single `generator` type, which can be recursive but doesn't implement symmetric transfer. Despite that, Folly users found the use of `Folly:::Generator` to be a lot more efficient than the eager algorithm they replaced with it.

`ranges-v3` also implements a `generator` type, which is never recursive and predates the work on move-only views and iterators [1], [2] which forces this implementation to ref-count the coroutine handler.

Our implementation [Implementation] consists of a single type that takes advantage of symmetric transfer to implement recursion - it notably works well with three different major standard libraries.

Performance & benchmarks

[*Note:* These benchmark results are fairly dated now - roughly a year old - and should be taken with a grain of salt. — *end note*]

Because implementations are still being perfected, and because performance is extremely dependant on whether HALO optimization (see P0981R0 [8]) occurs, it is difficult at this time to make definitive statements about the performance of the proposed design.

At the time of the writing of this paper, Clang is able to inline non-nested coroutines whether the implementation supports nested coroutines or not, while GCC never performs HALO optimization.

When the coroutine is not inlined, support for recursion does not noticeably impact performance. And, when the coroutine yields another generator, the performance of the recursive version is noticeably faster than yielding each element of the range. This is especially noticeable with deep recursion.

	Clang	Clang ST[1]	GCC	GCC ST[1]	MSVC	MSVC ST[1]
Single value	(1) 0.235	(2) 2.36	12.4	13.4	61.9	63.7
Single value, noinline (3)	13.5	13.7	14.1	15.2	63.8	64.4
Deep nesting	43670266.0	(4) 427955.0	58801348	338736	224052033	4760914

[1] `Symmetric transfer.`

The values are expressed in nanoseconds. However, please note that the comparison of the same result across compiler is not meaningful, notably because the MSVC results were obtained on different hardware. That being said, we observe:

- Only Clang can perform constant folding of values yielded by simple coroutine (1)
- When the `generator` supports symmetric transfer, clang is not able to fully inline the generator construction, but HALO is still performed (2).

- When HALO is not performed, the relative performance of both approaches is similar (3).
- Supporting recursion is greatly beneficial to nested/recursive algorithms (4).

The code for these benchmarks, as well as more detailed results, can be found on Github.

Wording

```
#define __cpp_lib_generator                    YYYYXXL // also in <generator>
```

Drafting Note: Modify [ranges.general] as follows:

❖ General [ranges.general]

This Clause describes components for dealing with ranges of elements.

The following subclauses describe range and view requirements, and components for range primitives and range generators as summarized in Table [tab:range.summary].

Drafting Note: Add a new row at the end of [tab:range.summary] "Range generators" with header `<generator>` referring to the new subclause [coroutine.generator] added below.

Drafting Note: Add the declaration of `ranges::elements_of` to the `<ranges>` synopsis:

❖ Header `<ranges>` synopsis [ranges.syn]

```
namespace std::ranges {
[...]

  template<input_or_output_iterator I, sentinel_for<I> S, subrange_kind K>
  inline constexpr bool enable_borrowed_range<subrange<I, S, K>> = true;

  // [range.dangling], dangling iterator handling
  struct dangling;

  // [elementsof.overview], class template elements_of
  template<range R, class Allocator = allocator<byte>>
  struct elements_of;

  template<range R>
  using borrowed_iterator_t = conditional_t<borrowed_range<R>, iterator_t<R>, dangling>;

  [...]
}
```

Drafting Note: Insert the following new subclause immediately after [range.dangling]:

❖ class template `elements_of` [ranges.elementsof]

❖ Overview [ranges.elementsof.overview]

Specializations of `elements_of` encapsulate a range and act as a tag in overload sets to disambiguate when a range should be treated as a sequence rather than a single value.

[*Example:*

```
std::generator<any> f(std::ranges::input_range auto&& rng) {
    co_yield rng; // yield rng as a single value
    co_yield std::ranges::elements_of(rng); // yield each element of rng
}
```

— *end example*]

```
namespace std::ranges {
  template<range R, class Allocator = allocator<byte>>
  struct elements_of {
    R range;
    [[no_unique_address]] Allocator allocator{};
  };

  template<class R, class Allocator = allocator<byte>>
  elements_of(R&&, Allocator = {}) -> elements_of<R&&, Allocator>;
}
```

Drafting Note: Add the following subclause to the end of [ranges]:

❖ Range Generators   [coroutine.generator]

❖ Overview   [coroutine.generator.overview]

`generator` presents a view of the elements yielded by the evaluation of a coroutine.

A `generator` generates a sequence of elements by repeatedly resuming the coroutine it was returned from. When the coroutine is resumed, it is executed until it reaches either a `co_-yield` statement or the end of the coroutine. Elements of the sequence are produced by the coroutine each time a `co_yield` statement is evaluated. When the `co_yield` statement is of the form `co_yield elements_of(rng)`, each element of the range `rng` is successively produced as an element of the generator.

[*Example:*

```
std::generator<int> iota(int start = 0) {
    while (true)
        co_yield start++;
}

void f() {
    for (auto i : iota() | std::views::take(3))
        std::cout << i << ' '; // prints 0 1 2
}
```

— *end example*]

❖ Header `<generator>` synopsis   [generator.syn]

20

```
#include <coroutine>
#include <ranges>

namespace std {
  // [coroutine.generator.class], class template generator
  template<class R, class V = void, class Allocator = void>
  class generator;

  template<class R, class V, class Allocator>
  constexpr bool ranges::enable_view<generator<R, V, Allocator>> = true;
}
```

�   Class template generator   [coroutine.generator.class]

```
namespace std {
  template<class R, class V = void, class Allocator = void>
  class generator {
    using value =          // exposition only
      conditional_t<is_void_v<V>, remove_cvref_t<R>, V>;
    using reference =    // exposition only
      conditional_t<is_void_v<V>, R&&, R>;
    class iterator;        // exposition only

  public:
    using yielded =
      conditional_t<is_reference_v<reference>, reference, const reference&>;

    class promise_type;

    generator(const generator&) = delete;
    generator(generator&& other) noexcept;

    ~generator();

    generator& operator=(const generator&) = delete;
    generator& operator=(generator&& other) noexcept;

    iterator begin();
    default_sentinel_t end() const noexcept;

  private:
    explicit generator(coroutine_handle<promise_type> coroutine) noexcept; // exposition only

    coroutine_handle<promise_type> coroutine_ = nullptr; // exposition only
  };
}
```

Mandates: `value` is a cv-unqualified object type.

Mandates: `reference` is either a reference type, or a cv-unqualified object type that models `copy_constructible`.

`Allocator` shall be `void`, or shall meet the `Cpp17Allocator` requirements.

Mandates: Let `RRef` denote `remove_reference_t<`*reference*`>&&` if *reference* is a reference type, or *reference* otherwise. Each of:

- `common_reference_with<`*reference*`&&, `*value*`&>`,
- `common_reference_with<`*reference*`&&, RRef&&>`, and
- `common_reference_with<RRef&&, const `*value*`&>`

is modeled. [*Note:* These requirements ensure the exposition-only *iterator* type can model `indirectly_readable` and thus `input_iterator`. — *end note*]

Specializations of `generator` model `view` and `input_range`.

The behavior of a program that adds a specialization for `generator` is undefined.

An instance of `generator` has an associated stack of coroutines, which is initially empty. A coroutine is associated with at most one `generator` instance at a given time.

�   Members   [generator.members]

```
explicit generator(coroutine_handle<promise_type> coro) noexcept;
```

> Initializes *coroutine_* with `coro`.

```
generator(generator&& other) noexcept;
```

> Initializes *coroutine_* with `exchange(other.`*coroutine_*`, {})`.

```
~generator();
```

> *Effects:* Equivalent to:
>
>  ```
>  if (coroutine_) {
>    coroutine_.destroy();
>  }
>  ```

```
generator& operator=(generator&& that) noexcept;
```

> *Effects:* Equivalent to:
>
>  ```
>  if (auto old = exchange(coroutine_, exchange(that.coroutine_, {}))) {
>    old.destroy();
>  }
>  ```
>
> *Returns:* `*this`.

```
iterator begin();
```

> *Preconditions:* *coroutine_* refers to a coroutine suspended at its initial suspend-point.
>
> *Effects:* Equivalent to:

```
      coroutine_.resume();
      return iterator(coroutine_);
```

> *Remarks:* This function pushes *coroutine_* onto the `generator`'s empty stack of associated coroutines.

> [*Note:* A program that calls `begin` more than once on the same generator has undefined behavior. *— end note*]

```
default_sentinel_t end() const noexcept;
```

> *Returns:* `default_sentinel`.

❖   class `generator::promise_type`   [coroutine.generator.promise]

```
template<class R, class V, class Allocator>
class generator<R, V, Allocator>::promise_type {
  friend generator;

  add_pointer_t<yielded> value_ = nullptr; // exposition only

public:
  generator get_return_object() noexcept;

  suspend_always initial_suspend() noexcept;

  auto final_suspend() noexcept;

  suspend_always yield_value(yielded value) noexcept;

  template<class Y = remove_reference_t<yielded>>
    requires is_rvalue_reference_v<yielded> &&
      constructible_from<remove_cvref_t<yielded>, const Y&>
        auto yield_value(const type_identity_t<Y>& lvalue);

  template<class R2, class V2, class Alloc2, class Unused>
    requires same_as<typename generator<T2, V2, Alloc2>::yielded, yielded>
      auto yield_value(ranges::elements_of<generator<T2, V2, Alloc2>&&, Unused> g) noexcept;

  template<ranges::input_range Rng, class Alloc2>
    requires convertible_to<ranges::range_reference_t<Rng>, yielded>
      auto yield_value(ranges::elements_of<Rng, Alloc2> r) noexcept;

  void await_transform() = delete;

  void return_void() noexcept {}

  void unhandled_exception();

  static void* operator new(size_t size)
    requires same_as<Allocator, void> || default_initializable<Allocator>;
```

23

```
  template<class Alloc, class... Args>
    requires same_as<Allocator, void> || convertible_to<Alloc, Allocator>
      static void* operator new(size_t size, allocator_arg_t, Alloc&& alloc, Args&...);

  template<class This, class Alloc, class... Args>
    requires same_as<Allocator, void> || convertible_to<Alloc, Allocator>
      static void* operator new(size_t size, This&, allocator_arg_t, Alloc&& alloc, Args&...);

  static void operator delete(void* pointer, size_t size) noexcept;
};
```

```
generator get_return_object() noexcept;
```

> *Returns:* `generator{coroutine_handle<promise_type>::from_promise(*this)}`.

```
suspend_always initial_suspend() noexcept;
```

> *Returns:* `{}`.

```
auto final_suspend() noexcept;
```

> *Preconditions:* The coroutine whose promise object is `*this` is at the top of the stack of associated coroutines of some `generator` instance x.

> *Returns:* An awaitable object of unspecified type whose member `await_suspend` removes the coroutine whose promise is `*this` from the top of x's stack of associated coroutines, and resumes execution of the new top-of-stack coroutine, if any.

```
suspend_always yield_value(yielded x) noexcept;
```

> *Effects:* Equivalent to: `value_ = addressof(x)`.

> *Returns:* `{}`.

```
template<class Y = remove_reference_t<yielded>>
  auto yield_value(const type_identity_t<Y>& lvalue)
    requires is_rvalue_reference_v<yielded> &&
      constructible_from<remove_cvref_t<yielded>, const Y&>;
```

> *Returns:* An object of an unspecified awaitable type ([expr.await]) that stores an object of type `remove_cv_t<Y>` direct-non-list-initialized with `lvalue`, whose member functions arrange for `value_` to point at that stored object before suspending the coroutine.

> *Throws:* Any exception thrown by the initialization of the stored object.

```
template<class T2, class V2, class Alloc2, class Unused>
  requires same_as<typename generator<T2, V2, Alloc2>::yielded, yielded>
    auto yield_value(ranges::elements_of<generator<T2, V2, Alloc2>&&, Unused> g) noexcept;
```

> *Preconditions:* The coroutine whose promise object is `*this` is at the top of the stack of associated coroutines of some `generator` instance x.

Let `r` be an lvalue that denotes the same object to which `g.range` refers.

Returns: An object of an unspecified awaitable type ([expr.await]) which takes ownership of the generator `r`, whose member `await_suspend` pushes `r.`*coroutine_* atop the stack of coroutines associated with `x` before resuming execution of `r.`*coroutine_*, and whose member `await_resume` rethrows any exception captured by a call to `r.`*coroutine_*`.promise()`'s member `unhandled_exception`.

[*Note:* Variables with automatic storage duration in the scope of the coroutine represented by `r.`*coroutine_* are destroyed before variables with automatic storage duration in the scope of the coroutine whose promise object is `*this`. — *end note*]

```
template<ranges::input_range Rng, class Alloc2>
  requires convertible_to<ranges::range_reference_t<Rng>, yielded>
    auto yield_value(ranges::elements_of<Rng, Alloc2> r) noexcept;
```

Effects: Equivalent to:

```
auto nested = [](allocator_arg_t, Alloc2, auto* range_ptr)
  -> generator<yielded, ranges::range_value_t<Rng>, Alloc2> {
    for (auto&& e : *range_ptr)
      co_yield static_cast<yielded>(std::forward<decltype(e)>(e));
  };
auto& rng = r.range;
return yield_value(ranges::elements_of(nested(
  allocator_arg, r.allocator, addressof(rng))));
```

```
void unhandled_exception();
```

Preconditions: The coroutine whose promise object is `*this` is at the top of the stack of associated coroutines of some `generator` instance `x`.

Effects: If the coroutine whose promise object is `*this` is the sole element of `x`'s stack of associated coroutines, equivalent to: `throw`. Otherwise, stores the result of `current_-exception()` where it can later be retrieved and rethrown by the `await_resume` member of the awaitable object returned from the `yield_value` call that pushed this coroutine onto `x`'s stack of associated coroutines.

```
static void* operator new(size_t size)
  requires same_as<Allocator, void> || default_initializable<Allocator>;
```

Let `A` be `allocator<void>` if `Allocator` denotes `void`, or `Allocator` otherwise. Let `BAlloc` be `allocator_traits<A>::template rebind_alloc<U>` where `U` denotes an unspecified type whose size and alignment are both `_STDCPP_DEFAULT_NEW_ALIGNMENT__`.

Effects: Initializes an allocator of type `BAlloc` with `A{}`, and uses that object to allocate the smallest number of blocks that provide sufficient storage for:

- a coroutine state of size `size`,
- if `allocator_traits<BAlloc>::is_always_equal::value` is `false`, space to store a copy of the allocator, and

25

- if `Allocator` denotes `void`, any additional state necessary to ensure that `operator delete` can later deallocate this memory block with an allocator equal to the allocator used here.

Returns: A pointer to the space allocated for the coroutine state.

```
template<class Alloc, class... Args>
  requires same_as<Allocator, void> || convertible_to<Alloc, Allocator>
    static void* operator new(size_t size, allocator_arg_t, Alloc&& alloc, Args&...);

template<class This, class Alloc, class... Args>
  requires same_as<Allocator, void> || convertible_to<Alloc, Allocator>
    static void* operator new(size_t size, This&, allocator_arg_t, Alloc&& alloc, Args&...);
```

Let `A` be `allocator<void>` if `Allocator` denotes `void`, or `Allocator` otherwise. Let `BAlloc` be `allocator_traits<A>::template rebind_alloc<U>` where `U` denotes an unspecified type whose size and alignment are both `_STDCPP_DEFAULT_NEW_ALIGNMENT__`.

Effects: Initializes an allocator of type `BAlloc` with `A(std::forward<Alloc>(alloc))`, and uses that object to allocate the smallest number of blocks that provide sufficient storage for:

- a coroutine state of size `size`,
- if `allocator_traits<BAlloc>::is_always_equal::value` is `false` or `default_initializable<BAlloc>` is `false`, space to store a copy of the allocator, and
- if `Allocator` denotes `void`, any additional state necessary to ensure that `operator delete` can later deallocate this memory block with an allocator equivalent to the allocator used here.

Returns: A pointer to the space allocated for the coroutine state.

```
static void operator delete(void* pointer, size_t size) noexcept;
```

Preconditions: `pointer` was returned from an invocation of one of the above overloads of `operator new` with a `size` argument equal to `size`.

Effects: Deallocates the block of allocator memory that includes the coroutine state denoted by `pointer` using an allocator equivalent to the one that was used to allocate it.

◆ Class template `generator::iterator`   [coroutine.generator.iterator]

```
template<class R, class V, class Allocator>
class generator<R, V, Allocator>::iterator {
public:
    using value_type = value;
    using difference_type = ptrdiff_t;

    iterator(iterator&& other) noexcept;

    iterator& operator=(iterator&& other) noexcept;
```

```
    reference operator*() const noexcept(is_nothrow_copy_constructible_v<reference>);

    iterator& operator++();
    void operator++(int);

    bool operator==(default_sentinel_t) const noexcept;

private:
    friend class generator;

    explicit iterator(coroutine_handle<promise_type> coroutine) noexcept; // exposition only

    coroutine_handle<promise_type> coroutine_; // exposition only
};
```

```
iterator(iterator&& other) noexcept;
```

> *Effects:* Initializes `coroutine_` with exchange(other.`coroutine_`, {}).

```
iterator& operator=(iterator&& other) noexcept;
```

> *Effects:* Equivalent to: `coroutine_` = exchange(other.`coroutine_`, {});

```
reference operator*() const noexcept(is_nothrow_copy_constructible_v<reference>);
```

> *Preconditions:* `coroutine_`.done() is false, and `coroutine_`

> Let p be the promise object of the coroutine at the top of the stack of coroutines associated with the `generator` whose stack of associated coroutines includes `coroutine_`.

> *Effects:* Equivalent to:

```
        return static_cast<reference>(*p.value_);
```

```
iterator& operator++();
```

> *Preconditions:* `coroutine_`.done() is false.

> *Effects:* Resumes the coroutine at the top of the stack of coroutines associated with the `generator` whose stack of associated coroutines includes `coroutine_`.

> *Returns:* return *this;

```
void operator++(int);
```

> *Preconditions:* `coroutine_`.done() is false.

> *Effects:* Equivalent to: ++*this.

```
bool operator==(default_sentinel_t) const noexcept;
```

> *Returns:* `coroutine_`.done().

```
explicit iterator(coroutine_handle<promise_type> coroutine) noexcept;
```

> *Effects:* Initializes *coroutine_* with coroutine.

References

[1] Casey Carter. P1456R1: Move-only views. https://wg21.link/p1456r1, 11 2019.

[2] Corentin Jabot. P1207R0: Movability of single-pass iterators. https://wg21.link/p1207r0, 8 2018.

[3] Corentin Jabot and Lewis Baker. P2168R3: generator: A synchronous coroutine generator compatible with ranges. https://wg21.link/p2168r3, 4 2021.

[4] Corentin Jabot, Eric Niebler, and Casey Carter. P1206R3: ranges::to: A function to convert any range to a container. https://wg21.link/p1206r3, 11 2020.

[5] Thomas Köppe. N4901: Working draft, standard for programming language c++. https://wg21.link/n4901, 10 2021.

[6] Gor Nishanov. P1681R0: Revisiting allocator model for coroutine lazy/task/generator. https://wg21.link/p1681r0, 6 2019.

[7] Barry Revzin and Tim Song. P2415R2: What is a `view`? https://wg21.link/p2415r2, 10 2021.

[8] Richard Smith and Gor Nishanov. P0981R0: Halo: coroutine heap allocation elision optimization: the joint response. https://wg21.link/p0981r0, 3 2018.

[9] Mathias Stearn. P2529R0: `generator` should have `T&& reference_type`. https://isocpp.org/files/papers/P2529R0.html, 1 2022.

[CppCoro] Lewis Baker. CppCoro: A library of C++ coroutine abstractions for the coroutines TS. https://github.com/lewissbaker/cppcoro

[Folly] Facebook. Folly: An open-source C++ library developed and used at Facebook. https://github.com/facebook/folly

[range] Eric Niebler. range-v3 Range library for C++14/17/20. https://github.com/ericniebler/range-v3

[Implementation] Casey Carter, Lewis Baker, Corentin Jabot. `std::generator` implementation. https://godbolt.org/z/qv9Yodhdx